## Math 114L

## Extra Credit Assignment

Spring 2010

1. Consider the model $\mathfrak{A}$ with $|\mathfrak{A}|=\mathbb{N},<^{\mathfrak{A}}=<$, and $+{ }^{\mathfrak{A}}=+$. Determine:
(a) For which variable assignments $s$ and which $d_{1}, d_{2}, d_{3}$ does

$$
\vDash_{\mathfrak{A}} v_{1}+v_{3}=v_{2}\left[s\left(v_{1} \mapsto d_{1}, v_{2} \mapsto d_{2}, v_{3} \mapsto d_{3}\right)\right] ?
$$

(b) For which variable assignments $s$ and which $d_{1}, d_{2}, d_{3}$ does

$$
\vDash_{\mathfrak{A}} \neg v_{1}+v_{3}=v_{2}\left[s\left(v_{1} \mapsto d_{1}, v_{2} \mapsto d_{2}, v_{3} \mapsto d_{3}\right)\right] ?
$$

(c) For which variable assignments $s$ and which $d_{1}, d_{2}$ does

$$
\vDash_{\mathfrak{A}} \forall v_{3} \neg v_{1}+v_{3}=v_{2}\left[s\left(v_{1} \mapsto d_{1}, v_{2} \mapsto d_{2}\right)\right] ?
$$

(d) For which variable assignments $s$ and which $d_{1}, d_{2}$ does

$$
\vDash_{\mathfrak{A}} \neg \forall v_{3} \neg v_{1}+v_{3}=v_{2}\left[s\left(v_{1} \mapsto d_{1}, v_{2} \mapsto d_{2}\right)\right] ?
$$

(e) For which variable assignments $s$ and which $d_{1}, d_{2}$ does

$$
\vDash_{\mathfrak{A}} \exists v_{3} v_{1}+v_{3}=v_{2}\left[s\left(v_{1} \mapsto d_{1}, v_{2} \mapsto d_{2}\right)\right] ?
$$

(f) For which variable assignments $s$ and which $d_{1}, d_{2}$ does

$$
\vDash_{\mathfrak{A}} v_{1}<v_{2}\left[s\left(v_{1} \mapsto d_{1}, v_{2} \mapsto d_{2}\right)\right] ?
$$

(g) For which variable assignments $s$ and which $d_{1}, d_{2}$ does

$$
\vDash_{\mathfrak{A}} v_{1}<v_{2} \rightarrow \exists v_{3} v_{1}+v_{3}=v_{2}\left[s\left(v_{1} \mapsto d_{1}, v_{2} \mapsto d_{2}\right)\right] ?
$$

(h) For which variable assignments $s$ and which $d_{1}$ does

$$
\vDash_{\mathfrak{A}} \forall v_{2}\left(v_{1}<v_{2} \rightarrow \exists v_{3} v_{1}+v_{3}=v_{2}\right)\left[s\left(v_{1} \mapsto d_{1}\right)\right] ?
$$

(i) For which variable assignments $s$ does

$$
\vDash_{\mathfrak{A}} \forall v_{1} \forall v_{2}\left(v_{1}<v_{2} \rightarrow \exists v_{3} v_{1}+v_{3}=v_{2}\right)[s] ?
$$

(j) Does

$$
\vDash_{\mathfrak{A}} \forall v_{1} \forall v_{2}\left(v_{1}<v_{2} \rightarrow \exists v_{3} v_{1}+v_{3}=v_{2}\right) ?
$$

2. Does $\vDash_{\mathfrak{A}} \exists v_{1} \forall v_{2} v_{2}<v_{2}+v_{1}$ ? Break the question into steps, as in the previous part.
3. Let $\alpha$ be the formula $\forall x \forall y x<y \rightarrow \exists z x<z \wedge z<y$, let $\beta$ be the formula $\forall x \exists y x<y$, and let $\gamma$ be the formula $\forall x \forall y x<y \rightarrow \neg x=y$. Show that:
(a) $\alpha ; \gamma \nvdash \beta$
(b) $\alpha ; \gamma \nvdash \neg \beta$
(c) $\beta ; \gamma \nvdash \alpha$
(d) $\beta ; \gamma \nvdash \neg \alpha$
4. Write down a deduction showing that $\forall x \forall y \phi \vdash \forall x \phi$. (Don't just prove that the deduction exists; actually write it down as a list of formulas.)
5. Write down a deduction showing that $\forall x \forall y \phi \vdash \forall y \forall x \phi$. (Hint: the fact that there is such a deduction follows from the previous part and the generalization theorem. Use the proof of the generalization theorem to transform your deduction from the previous part, step by step, into the deduction for this part.)
6. Write down a deduction showing that $\vdash \forall x \forall y \phi \rightarrow \forall y \forall x \phi$. (Hint: similar to the last part, but with the deduction theorem.)
7. Prove, carefully, that if $\Gamma$ is a consistent set of first-order formulas in a language $\mathcal{L}$ (and $\mathcal{L}$ has only countably many predicate and function symbols) then there is a consistent set $\Delta \supseteq \Gamma$ such that for every formula $\phi$, either $\phi \in \Delta$ or $\neg \phi \in \Delta$.
