



\* is a short deduction from axioms similar to the other two, which I'm omitting for space.  
 For convenience, there are some named deductions here. In particular,  $d(x + 1)$  is a deduction of

$$\Lambda, \exists y \forall w \neg A(x, y, w), \forall y \exists z A(x + 1, y, z),$$

with  $x$  understood as a parameter other than 0. By  $d(3)$ , for instance, we mean the same deduction with every instance of  $x$  replaced by 3. (By  $d(0)$  we mean the deduction of  $\Lambda, \forall y \exists z A(0, y, z)$ .)

Now suppose we plug in some not-completely-trivial values of  $x$  and  $y$ , say  $x = 3, y = 2$ , and apply cut-elimination. We start by unraveling the induction; depending on our exact algorithm, we get something like the following.

We end up with a deduction of  $\Lambda, \forall y, \exists z A(2, y, z)$ :

Unraveling the induction, we get a deduction like the following (depending on our exact algorithm):

$$\begin{array}{c} \vdots \\ d(0) \frac{\Lambda, \forall y \exists z A(0, y, z)}{\text{Cut}} \quad d(1) \frac{\Lambda, \exists y \forall w \neg A(0, y, w), \forall y \exists z A(1, y, z)}{\text{Cut}} \quad d(2) \frac{\Lambda, \exists y \forall w \neg A(1, y, w), \forall y \exists z A(2, y, z)}{\text{Cut}} \\ \vdots \\ \Lambda, \forall y \exists z A(2, y, z) \end{array}$$

We'll call this  $p(2)$ . More interestingly, we get the following deduction:

$$\begin{array}{c} \frac{A(2, 1, z), \neg A(2, 1, z)}{\Lambda, \neg A(2, 1, z), A(3, 0, z)} \quad \frac{\neg A(3, 0, z), A(3, 0, z)}{\Lambda, \neg A(2, z, w), \neg A(3, 0, z), \exists z A(3, 1, z)} \quad \frac{*}{\Lambda, \neg A(2, z, w), \neg A(3, 1, z), A(3, 2, w)} \\ \frac{\Lambda, \neg A(2, 1, z), A(3, 0, z)}{\Lambda, \neg A(2, 1, z), \exists z A(3, 0, z)} \quad \frac{\Lambda, \neg A(2, z, w), \neg A(3, 0, z), \exists z A(3, 1, z)}{\Lambda, \forall w \neg A(2, z, w), \neg A(3, 0, z), \exists z A(3, 1, z)} \quad \frac{\Lambda, \neg A(2, z, w), \neg A(3, 1, z), \exists z A(3, 2, z)}{\Lambda, \forall w \neg A(2, z, w), \neg A(3, 1, z), \exists z A(3, 2, z)} \\ \frac{\Lambda, \forall z \neg A(2, 1, z), \exists z A(3, 0, z)}{\Lambda, \exists y \forall z \neg A(2, y, z), \exists z A(3, 0, z)} \quad \frac{\Lambda, \exists y \forall w \neg A(2, y, w), \neg A(3, 0, z), \exists z A(3, 1, z)}{\Lambda, \exists y \forall w \neg A(2, y, w), \forall z \neg A(3, 0, z), \exists z A(3, 1, z)} \quad \frac{\Lambda, \forall w \neg A(2, z, w), \neg A(3, 1, z), \exists z A(3, 2, z)}{\Lambda, \exists y \forall w \neg A(2, y, w), \neg A(3, 1, z), \exists z A(3, 2, z)} \\ \text{Cut} \frac{\Lambda, \exists y \forall z \neg A(2, y, z), \exists z A(3, 0, z)}{\text{Cut}} \quad \frac{\Lambda, \exists y \forall w \neg A(2, y, w), \forall z \neg A(3, 0, z), \exists z A(3, 1, z)}{\text{Cut}} \quad \frac{\Lambda, \exists y \forall w \neg A(2, y, w), \forall z \neg A(3, 1, z), \exists z A(3, 2, z)}{\text{Cut}} \end{array}$$

We'll call this  $a$ . Our unraveled deduction is the result of cutting these together:

$$\begin{array}{c} \vdots \\ p(2) \frac{\Lambda, \forall y \exists z A(2, y, z)}{\text{Cut}} \quad a \frac{\Lambda, \exists y \forall w \neg A(2, y, w), \exists z A(3, 2, z)}{\text{Cut}} \\ \vdots \\ \Lambda, \exists z A(3, 2, z) \end{array}$$

The cut-elimination procedure says we should follow the existential side of the cuts up to an actual  $\exists$  introduction rule; in this case it's the inference rule in  $a$  marked with  $\dagger$ . We permute the introduction rule down to match its cut, which in this case is all the way to the bottom.



$$\begin{array}{c}
\dots \\
\frac{\neg A(2, 5, w), A(2, 5, w)}{\Lambda, A(3, 0, 5) \wedge A(2, 5, w), \neg A(2, 5, w)} \quad \frac{\Lambda, A(3, 0, 5)}{\neg A(3, 1, w), A(3, 1, w)} \\
\frac{\Lambda, A(3, 0, 5) \wedge A(2, 5, w) \wedge \neg A(3, 1, w), \neg A(2, 5, w), A(3, 1, w)}{\Lambda, \neg A(2, 5, w), A(3, 1, w)} \quad \frac{\Lambda, \neg A(2, z, w), \neg A(3, 1, z), A(3, 2, w)}{\Lambda, \neg A(2, z, w), \neg A(3, 1, z), \exists z A(3, 2, z)} \\
\frac{\Lambda, \neg A(2, 5, w), \exists z A(3, 1, z)}{\Lambda, \forall w \neg A(2, 5, w), \exists z A(3, 1, z)} \quad \frac{\Lambda, \forall w \neg A(2, z, w), \neg A(3, 1, z), \exists z A(3, 2, z)}{\Lambda, \exists y \forall w \neg A(2, y, w), \neg A(3, 1, z), \exists z A(3, 2, z)} \\
\frac{\Lambda, \exists y \forall w \neg A(2, y, w), \exists z A(3, 1, z)}{\Lambda, \exists y \forall w \neg A(2, y, w), \forall z \neg A(3, 1, z), \exists z A(3, 2, z)} \\
\frac{\Lambda, \forall y \exists z A(2, y, z)}{\Lambda, \exists z A(3, 2, z)} \quad \text{Cut} \quad \frac{\Lambda, \exists y \forall w \neg A(2, y, w), \exists z A(3, 2, z)}{\Lambda, \exists z A(3, 2, z)}
\end{array}$$

The next cut to be eliminated is the one corresponding the existential quantifier at †, which is being cut with the copy of  $p(2)$ . Something exciting has finally happened, so we should take note: the value for  $y$  we're about to substitute into  $p(2)$  is the value 5, which is itself the value of the Ackermann function  $A(2, 1) = 5$ . If we think of  $p(2)$  as a function  $y \mapsto A(2, y)$ , we're now iterating this function.

Once we start iterating functions—or equivalently, eliminating the same cut we already eliminated ones—we should be worried about non-termination. But remember what happened when we eliminated the cut over  $\forall y \exists z A(2, y, z)$  the first time: we created a new copy of the same cut, but *with a lower ordinal*. For thoroughness, let's do this one more time; we permute the  $\exists$  introduction down:

$$\begin{array}{c}
\dots \\
\frac{\neg A(2, 5, w), A(2, 5, w)}{\Lambda, A(3, 0, 5) \wedge A(2, 5, w), \neg A(2, 5, w)} \quad \frac{\Lambda, A(3, 0, 5)}{\neg A(3, 1, w), A(3, 1, w)} \\
\frac{\Lambda, A(3, 0, 5) \wedge A(2, 5, w) \wedge \neg A(3, 1, w), \neg A(2, 5, w), A(3, 1, w)}{\Lambda, \neg A(2, 5, w), A(3, 1, w)} \quad \frac{\Lambda, \neg A(2, z, w), \neg A(3, 1, z), A(3, 2, w)}{\Lambda, \neg A(2, z, w), \neg A(3, 1, z), \exists z A(3, 2, z)} \\
\frac{\Lambda, \neg A(2, 5, w), \exists z A(3, 1, z)}{\Lambda, \forall w \neg A(2, 5, w), \exists z A(3, 1, z)} \quad \frac{\Lambda, \forall w \neg A(2, z, w), \neg A(3, 1, z), \exists z A(3, 2, z)}{\Lambda, \exists y \forall w \neg A(2, y, w), \neg A(3, 1, z), \exists z A(3, 2, z)} \\
\frac{\Lambda, \exists y \forall w \neg A(2, y, w), \exists z A(3, 1, z)}{\Lambda, \exists y \forall w \neg A(2, y, w), \forall z \neg A(3, 1, z), \exists z A(3, 2, z)} \\
\frac{\Lambda, \forall y \exists z A(2, y, z)}{\Lambda, \exists z A(3, 2, z)} \quad \text{Cut} \quad \frac{\Lambda, \forall w \neg A(2, 5, w), \exists y \forall w \neg A(2, y, w), \exists z A(3, 2, z)}{\Lambda, \exists y \forall w \neg A(2, y, w), \exists z A(3, 2, z)}
\end{array}$$

and then eliminate the cut:

$$\begin{array}{c}
\vdots \\
\frac{\neg A(2, 5, w), A(2, 5, w)}{\Lambda, A(3, 0, 5) \wedge A(2, 5, w), \neg A(2, 5, w)} \quad \frac{\vdots}{\Lambda, A(3, 0, 5)} \\
\frac{\Lambda, A(3, 0, 5) \wedge A(2, 5, w), \neg A(2, 5, w)}{\Lambda, A(3, 0, 5) \wedge A(2, 5, w) \wedge \neg A(3, 1, w), \neg A(2, 5, w), A(3, 1, w)} \quad \frac{\neg A(3, 1, w), A(3, 1, w)}{\Lambda, \neg A(2, 5, w), A(3, 1, w)} \\
\frac{\Lambda, \neg A(2, 5, w), A(3, 1, w)}{\Lambda, \neg A(2, 5, w), \exists z A(3, 1, z)} \\
\frac{\Lambda, \neg A(2, 5, w), \exists z A(3, 1, z)}{\Lambda, \forall w \neg A(2, 5, w), \exists z A(3, 1, z)} \quad \text{Cut} \\
\frac{\Lambda, \forall w \neg A(2, 5, w), \exists z A(3, 1, z)}{\Lambda, \forall w \neg A(2, 5, w), \exists y \forall w \neg A(2, y, w), \exists z A(3, 2, z)} \\
\frac{\Lambda, \forall w \neg A(2, 5, w), \exists z A(3, 2, z)}{\Lambda, \exists z A(3, 2, z)} \\
\frac{\Lambda, \exists z A(3, 2, z)}{\Lambda, \exists z A(2, 5, z)} \\
\frac{\Lambda, \exists z A(2, 5, z)}{\Lambda, \exists z A(2, 5, z)} \quad p(2, 5) \\
\frac{\Lambda, \forall y \exists z A(2, y, z)}{\Lambda, \forall y \exists z A(2, y, z)} \quad p(2) \\
\frac{\Lambda, \forall y \exists z A(2, y, z)}{\Lambda, \exists z A(3, 2, z)}
\end{array}$$

Again, we've made a new copy of  $p(2)$ , meaning we'll get to plug in a third value for  $y$ , but we've pushed the cut further up in the proof (by eliminating the  $\exists$  introduction down), meaning it has a smaller ordinal, and therefore we've made progress.

Going back to the beginning of our construction, let's consider what happened. We think of  $p(2)$ , the proof of  $\Lambda, \forall y \exists z A(2, y, z)$ , as corresponding to having the function  $y \mapsto A(2, y)$ . We cut it with a proof of  $\Lambda, \exists y \forall w \neg A(2, y, w), \exists z A(3, n, z)$  for some  $n$ . A single lower-rank induction axiom over  $n$  means that we introduced  $\exists y \forall w \neg A(2, y, w)$   $n$  separate times (phrased differently, the proof of  $\Lambda, \exists y \forall w \neg A(2, y, w), \forall y \exists z A(3, y, z)$  morally contains  $\omega$  introductions of  $\exists y \forall w \neg A(2, y, w)$ ; note that we do mean the ordinal  $\omega$ , not the cardinal  $\aleph_0$ —the introductions are still placed in a well-founded ordering). We then have to eliminate the cut over  $\forall y \exists z A(2, y, z)$   $n$  times, which corresponds to iteration:  $A(2, A(2, \dots))$ ,  $n$  times.